

What's the Story in EBS Glory: Evolutions and Lessons in Building Cloud Block Store

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Background: Elastic Block Store

● EBS

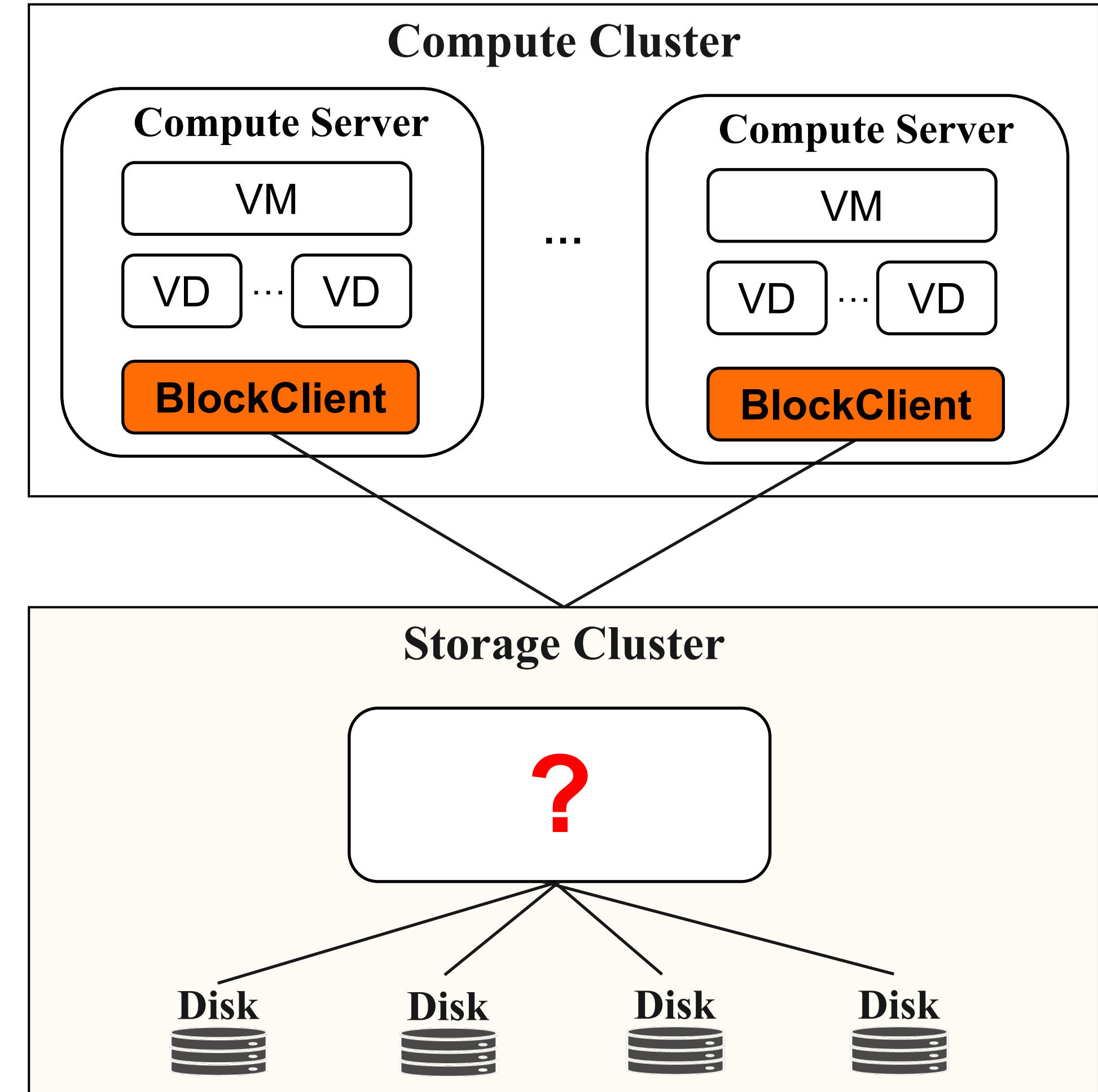
- ✓ VM: Virtual Machine
- ✓ VD: Virtual Disk

● Goal

- ✓ High Performance
- ✓ High Elasticity
- ✓ High Availability

● Compute-Storage Disaggregation

- ✓ **VMs** and **VDs** are on different clusters



Evolutions of EBS

Elasticity: A Tale of Four Metrics

Other Topics

EBS1: an Initial Foray

● Design Goals

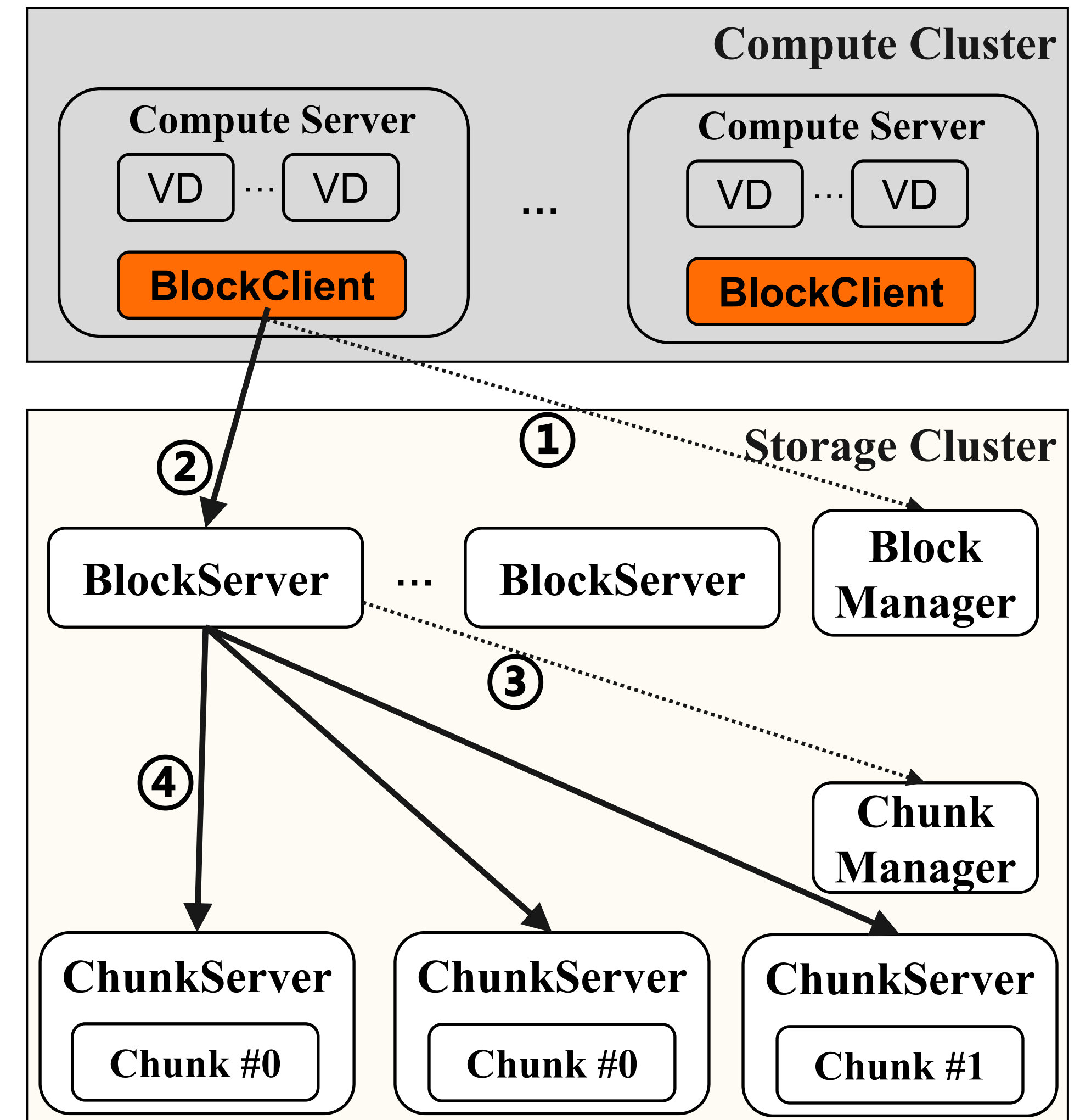
- ✓ **Straightforward** design for fast development/deployment

● Architecture

- ✓ VD space is partitioned into fixed-size **Chunks** (64 MiB)
- ✓ Two-layer: Blockserver + Chunkserver
- ✓ Each Chunk is an **Ext4 file**

● Features

- ✓ **In-place** updates: VD = Ext4 files
- ✓ **N(VDs)-to-1**(blockserver) binding



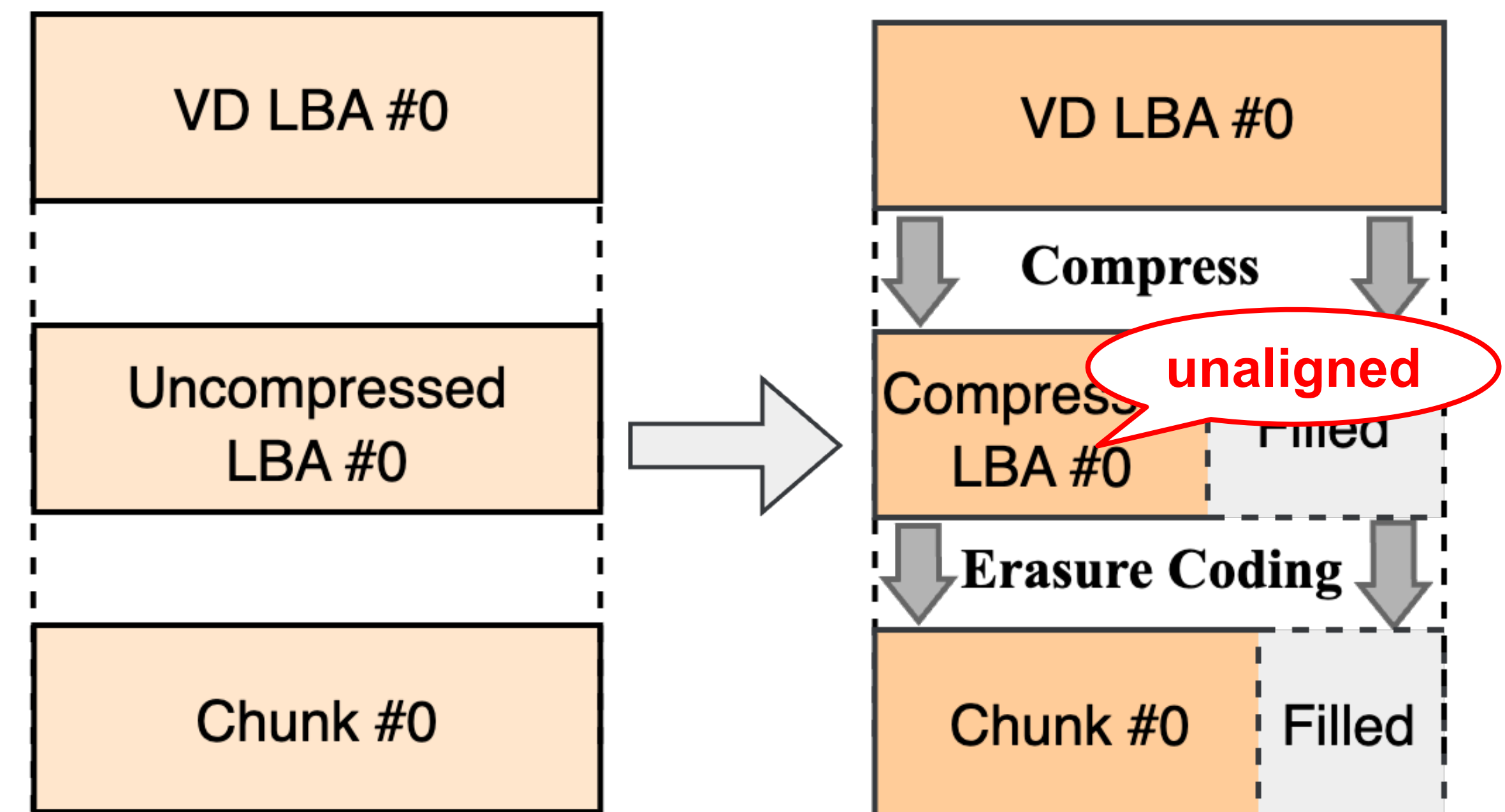
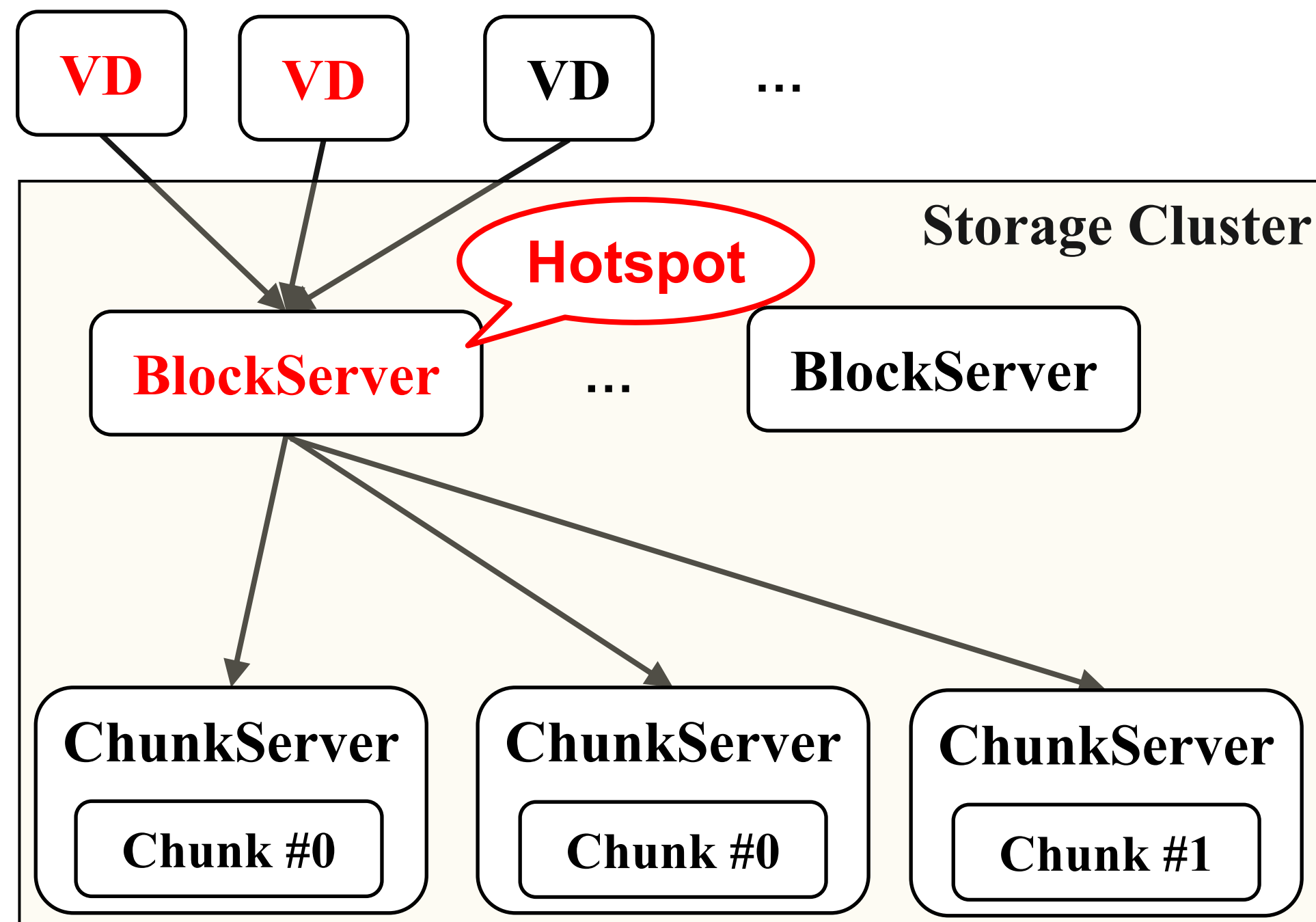
EBS1: An Initial Foray

● Deployment

- ✓ Released in **2012**, served over **1 million** VD's and stored hundreds of PBs of data across **hundreds of clusters**

● Limitations

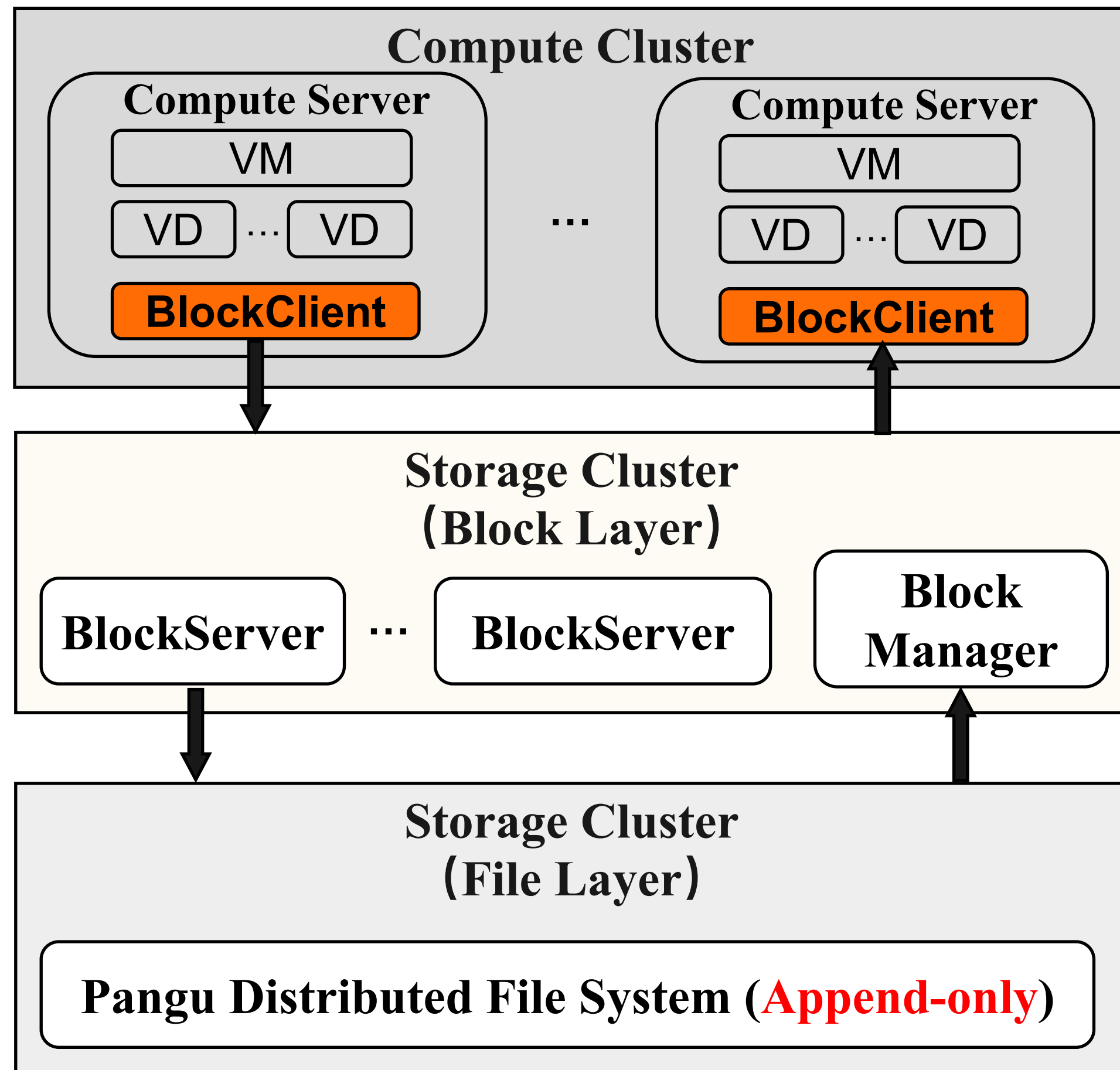
- ✓ **N-to-1 mapping** leads to **a single hot-point** bottlenecks and restricts performance
- ✓ **In-place updates** hinder the implementation of **compression and EC**, thereby reducing cost-efficiency



EBS2: Speedup with Space Efficiency

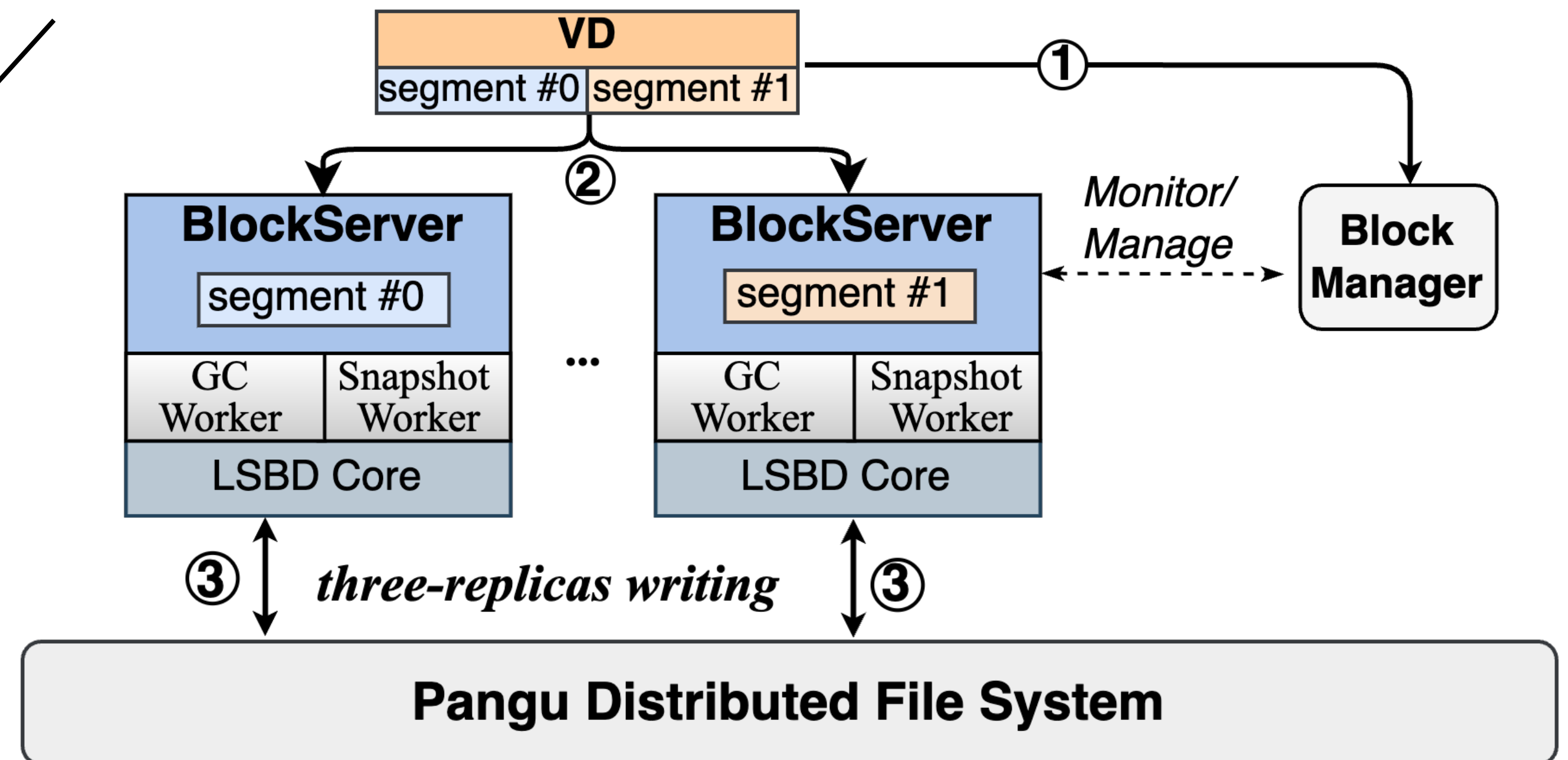
● Design Goals

- ✓ High **performance** and high **space efficiency**



● Key Designs

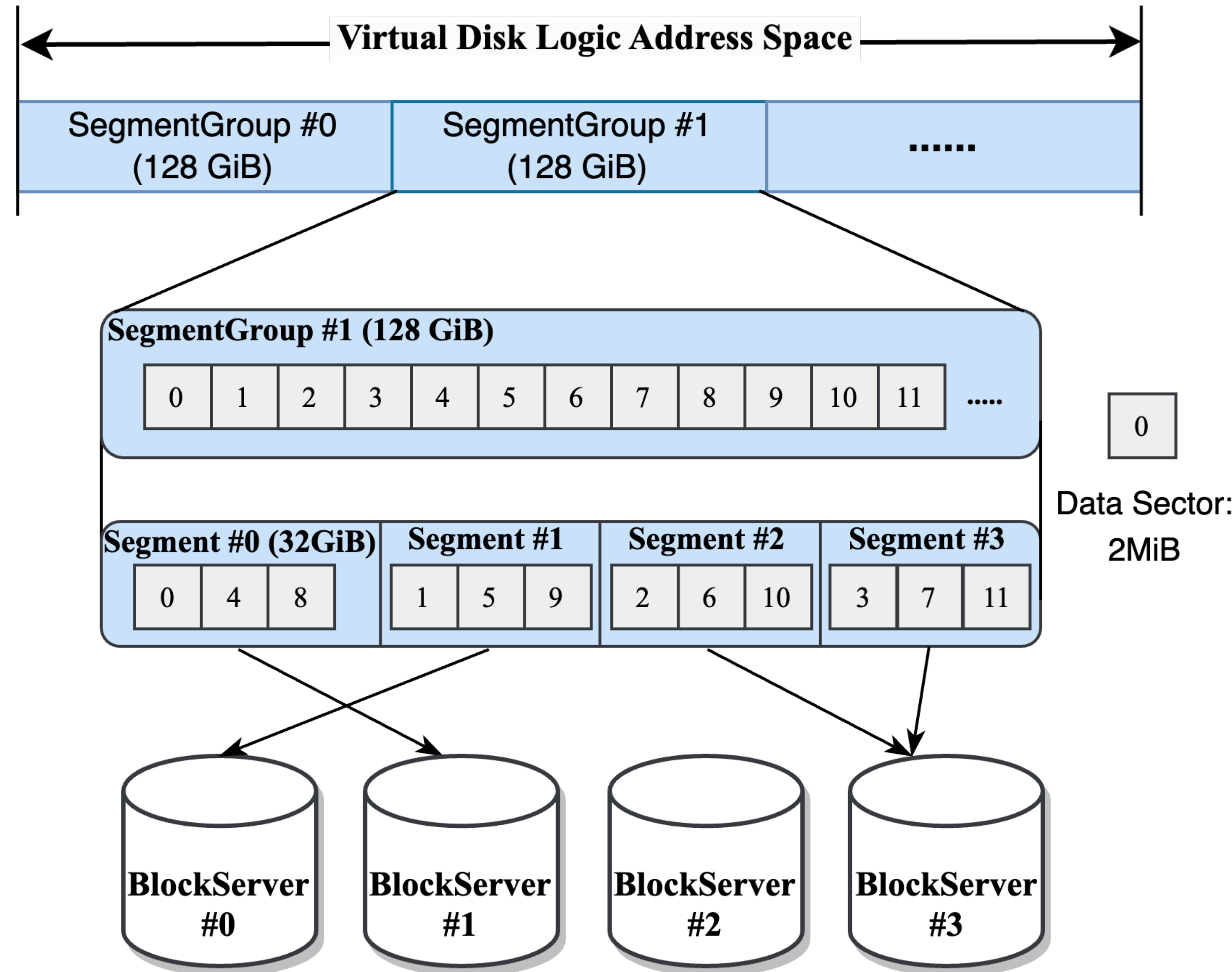
- ✓ **Disk segmentation**
- ✓ **Log-structured Block Device (LSBD)**
- ✓ **GC with EC/Compression**



EBS2: Speedup with Space Efficiency

● Disk Segmentation

- ✓ The entire VD logic space is divided into multiple contiguous **SegmentGroups**
- ✓ Each **SegmentGroup** is organized as a series of **Data Sectors**
- ✓ Data Sectors are allocated to the **Segments** in a Round-Robin Fashion
- ✓ BlockServers operate at the granularity of **Segments**

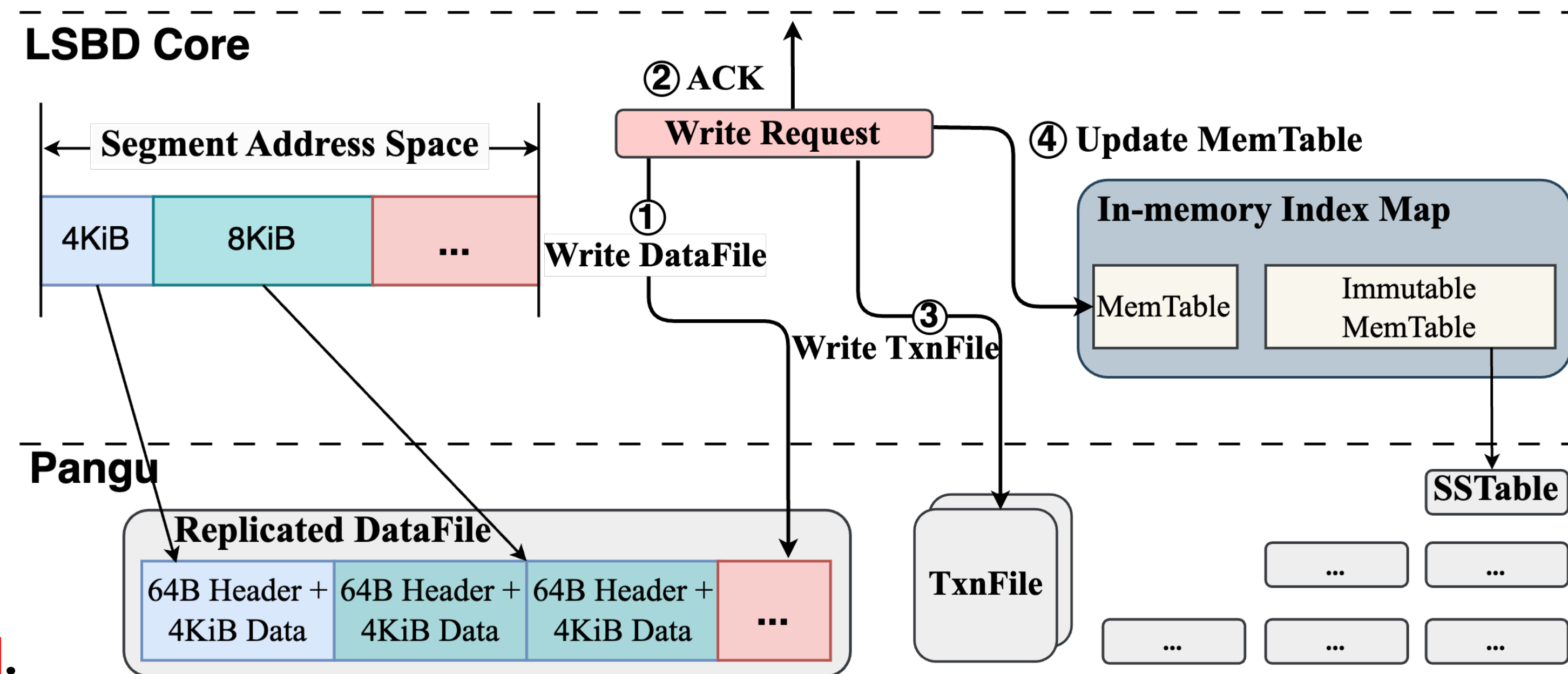


● Log-structured Block Device

✓ DataFile = (4KB data + 64B Header) x N

✓ Txnfile for speeding up failover

✓ In-memory Index Map for speeding up **read**.

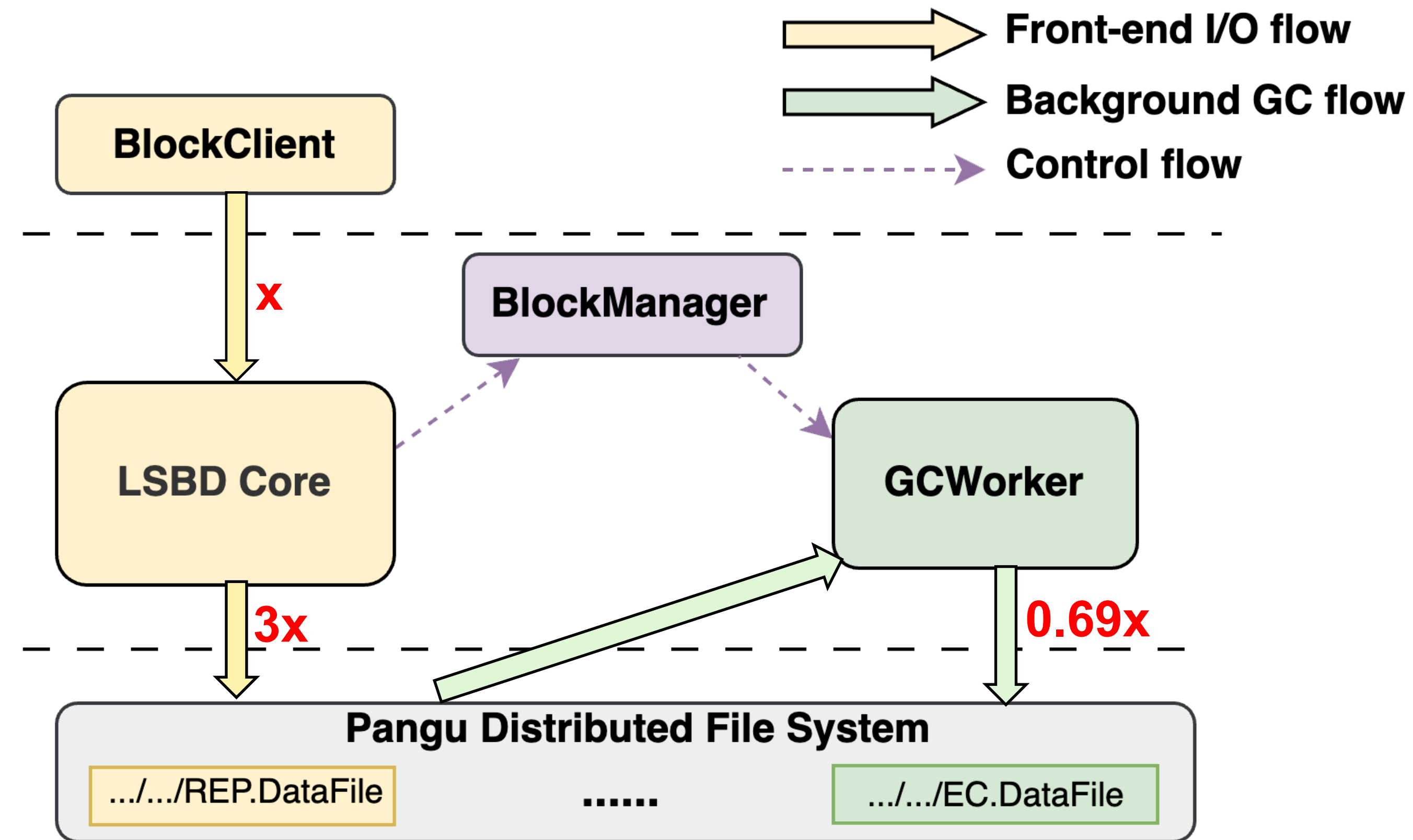


EBS2: Speedup with Space Efficiency

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GC with EC/Compression

- ✓ LSBD splits traffic into **frontend** (i.e., client I/Os) and **backend** (i.e., GC and compression)
- ✓ GC runs at the granularity of **DataFiles**
- ✓ GC converts the “REP.DataFiles” to “EC.DataFiles” with **EC(8, 3)** and **LZ4/ZSTD** compression algorithms



$$SpaceCost_{EBS1} = 3$$

$$SpaceCost_{EBS2} = 1(\text{original}) \times 0.5(\text{compressed}) \times \frac{8+3}{8} (\text{EC}) = 0.69$$

EBS2: Speedup with Space Efficiency

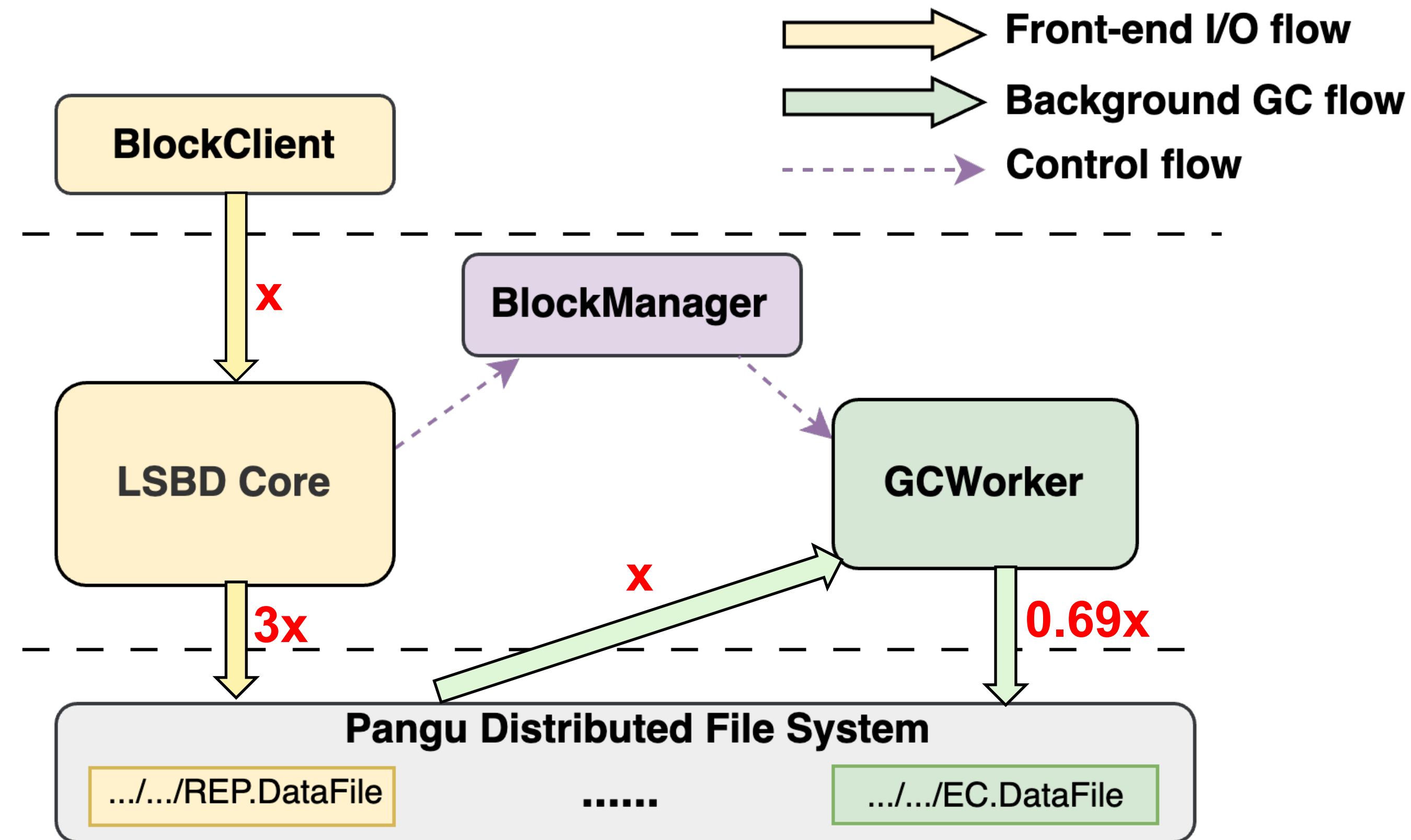
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Deployment

- ✓ **100μs** avg. write latency and **1 million** IOPS per VD.
- ✓ Over **500** clusters and served for **2 million** VDs.
- ✓ Low to **1.29** data replicas.

Limitations

- ✓ Traffic amplification up to **4.69**.
- ✓ As the cost per GiB of SSD decreases, cloud storage has shifted from **space-sensitive** to **traffic-sensitive**.



$$TrafficAmplification_{EBS1} = 3x \div x = 3$$

$$TrafficAmplification_{EBS2} = (3x + x + 0.69x) \div x = 4.69$$

EBS2 with Foreground EC/Compression?

- **Fragmented requests prevent Online Compress-EC**
 - ✓ EC requires the raw data blocks to typically be at least **16KB**
 - ✓ Nearly **70%** of write requests are smaller than 16KB
 - ✓ Waiting for merging incurs **extra latency (ranging from 10us to 100ms)**
- **CPU-based compression is slow**
 - ✓ 16KB-sized data blocks compression = **25us** for CPUs
 - ✓ **CPU resource** contention leads to lower throughput

EBS3: Foreground EC/Compression

● Design Goals

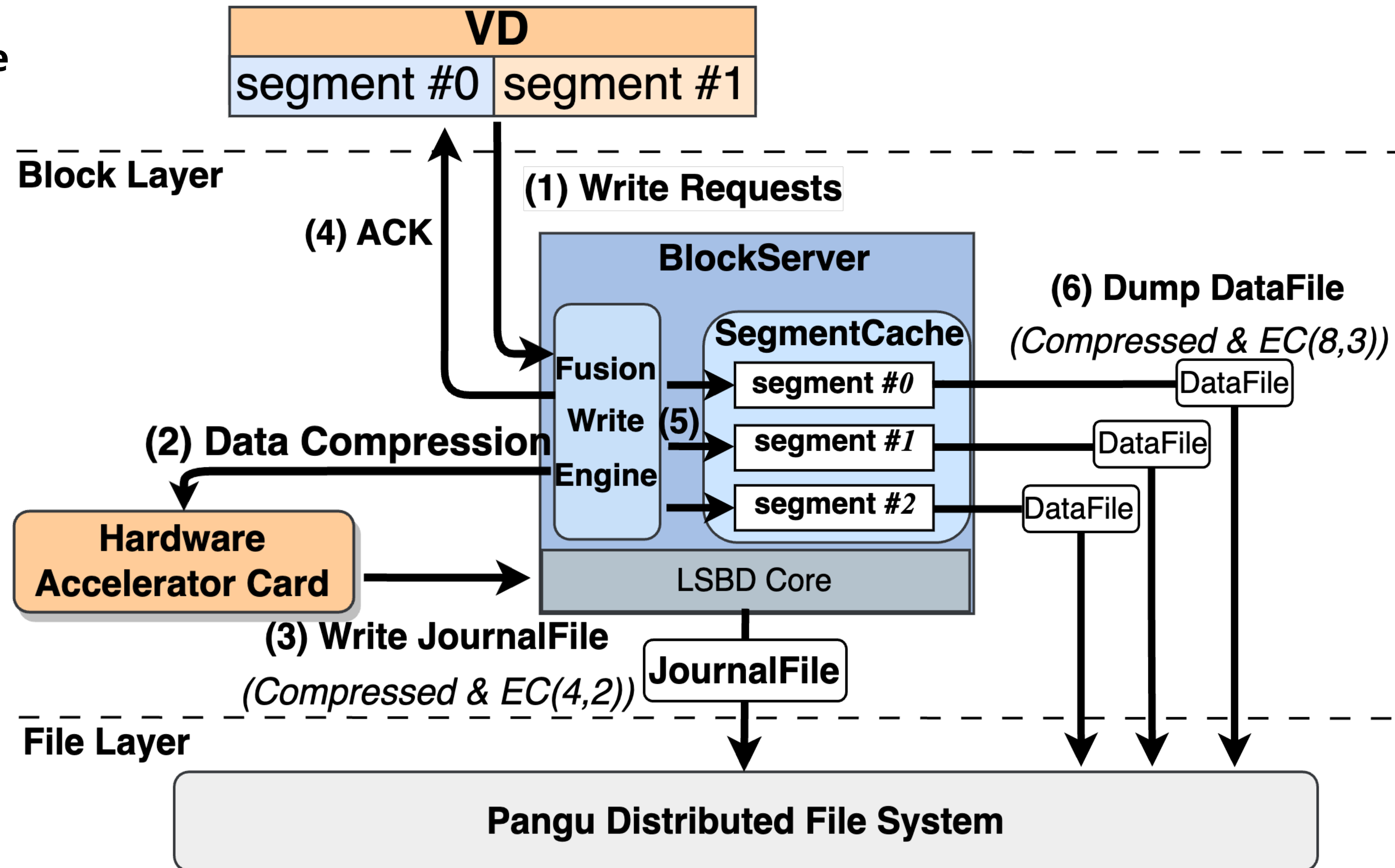
- ✓ Lower traffic consumption and storage space costs
- ✓ No performance loss

● Key Designs

- ✓ Bifurcated write path
- ✓ Fusion Write Engine
- ✓ FPGA-based compression offloading

● Deployment

- ✓ Over **100** clusters for **500K** VDs
- ✓ Data replicas reduced to **0.77**



EBS3: Foreground EC/Compression

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● Design Goals

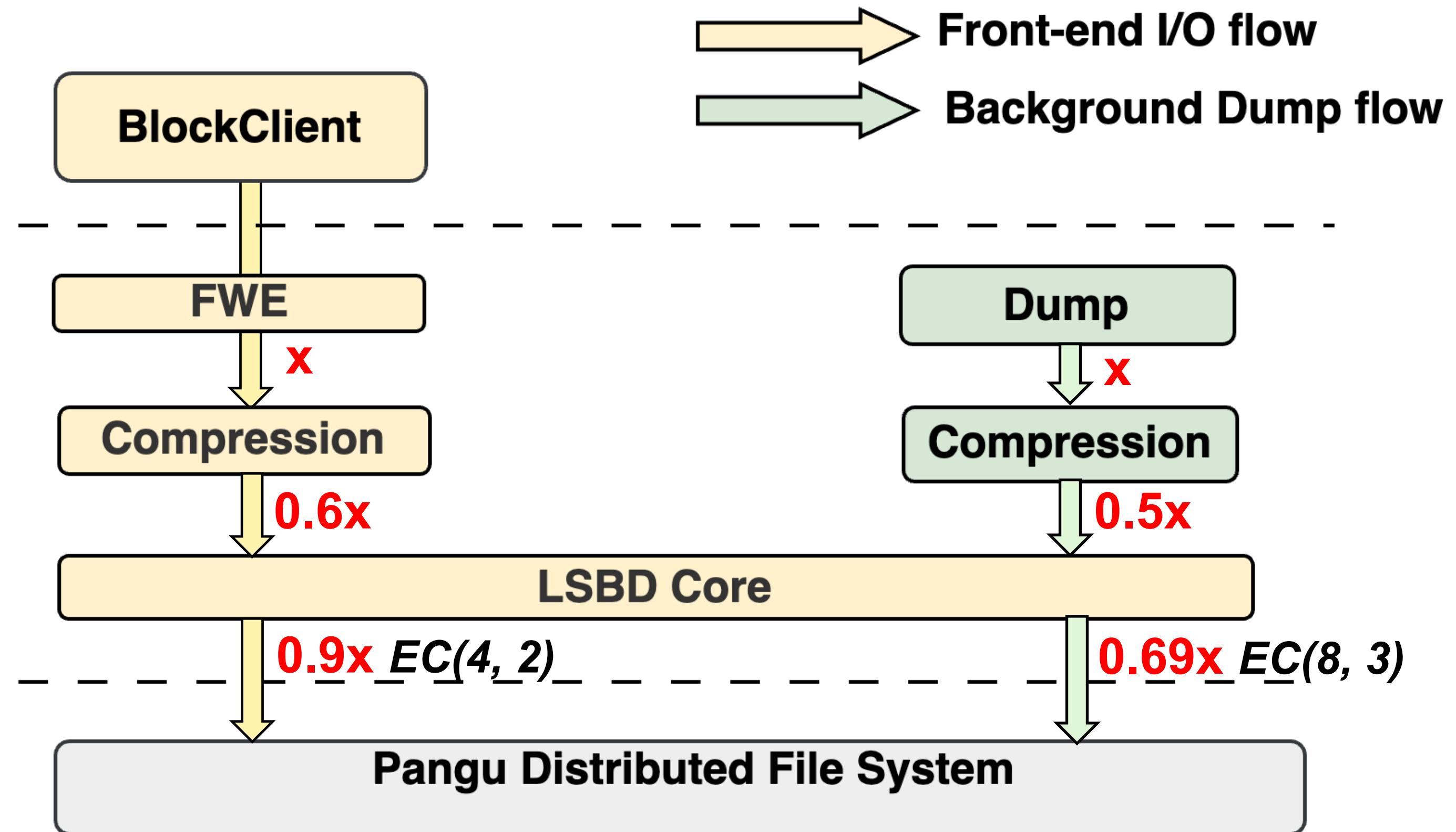
- ✓ Lower traffic consumption and storage space costs
- ✓ No performance loss

● Key Designs

- ✓ Fusion Write Engine
- ✓ FPGA-based compression offloading
- ✓ Traffic reduced from 4.69 to 1.59

● Deployment

- ✓ Over 100 clusters for 500K VD's
- ✓ Data replicas reduced to 0.77



$$TrafficAmplification_{EBS2} = (3x + x + 0.69x) \div x = 4.69$$

$$TrafficAmplification_{EBS3} = (0.9x + 0.69x) \div x = 1.59$$

Comparison of Three Generations of EBS

	EBS1	EBS2	EBS3
Avg. Latency	Millisecond Level	Hundred-microsecond Level	Hundred-microsecond Level
MAX. IOPS / Throughput	25,000	1,000,000	1,000,000
Key Features	In-place updates N-to-1mapping	Background EC & Compression	Foreground EC & Compression
Space Cost (Replicas per Data)	3	1.29	0.77
Traffic Amplification	3	4.69	1.59

Evolving Journey of EBS

Elasticity: A Tale of Four Metrics

Other Topics

Metrics #1: Latency

- **Elasticity of latency is coarse-grained**

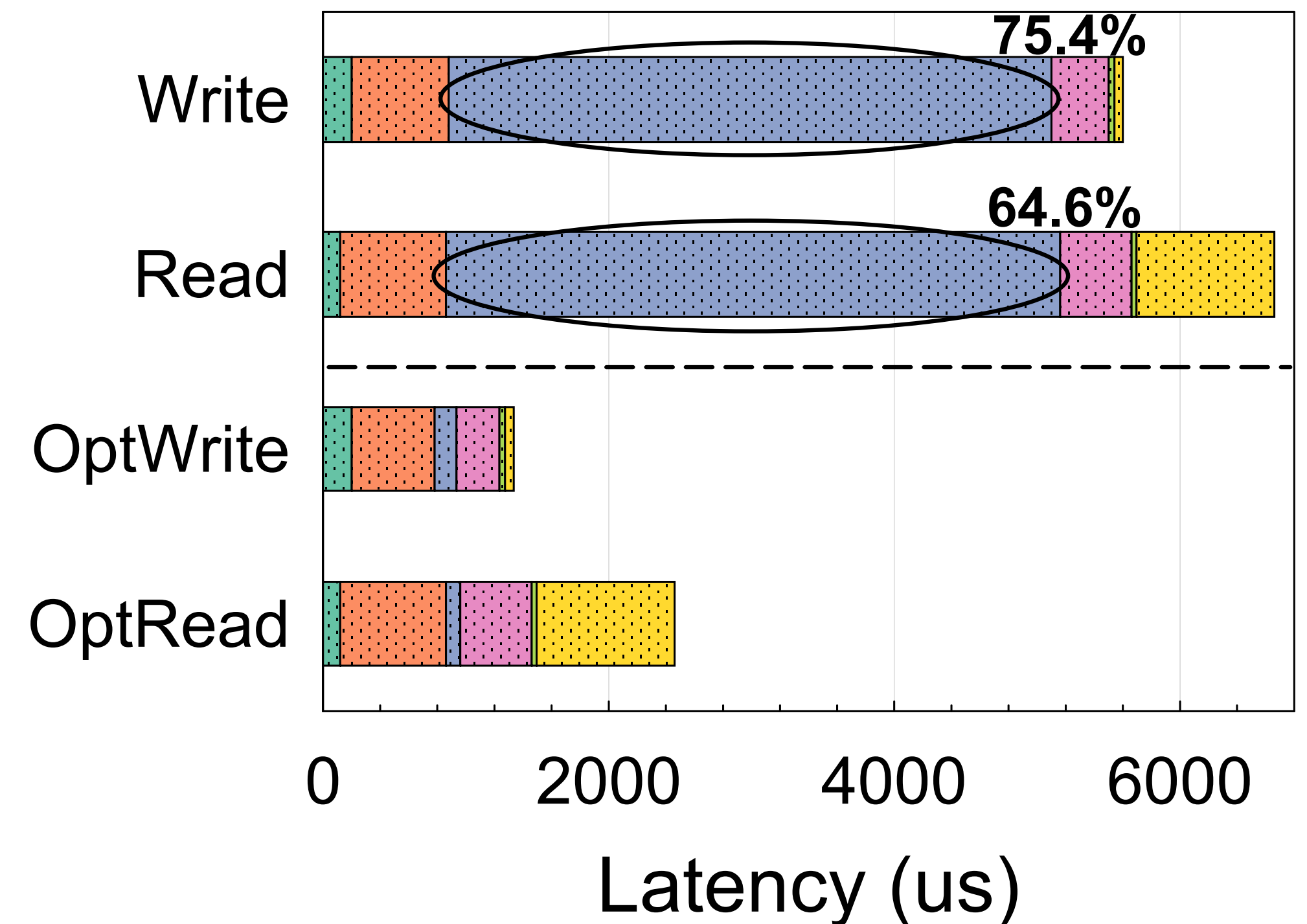
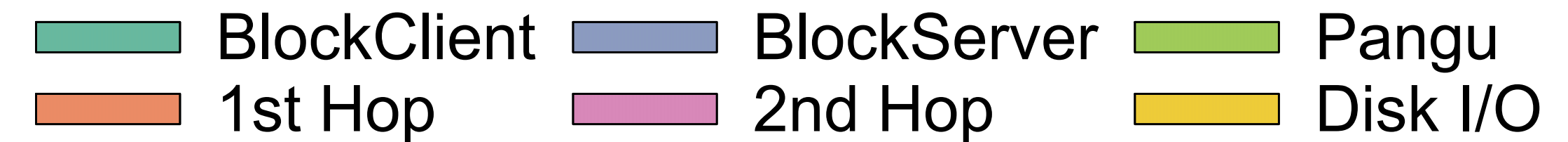
- ✓ Defined by the architectures

- **EBSX**

- ✓ **Shorten the path** (*e.g., skip a network hop*)
- ✓ **Use faster devices** (*e.g., PMem instead of SSD*)
- ✓ **Simple and efficient** data consistency protocol

- **Tail latency**

- ✓ **Software-induced** tail latency can be the dominant
- ✓ Separate client IOs from background tasks (*e.g., GC*)



99.999th Percentile Latency

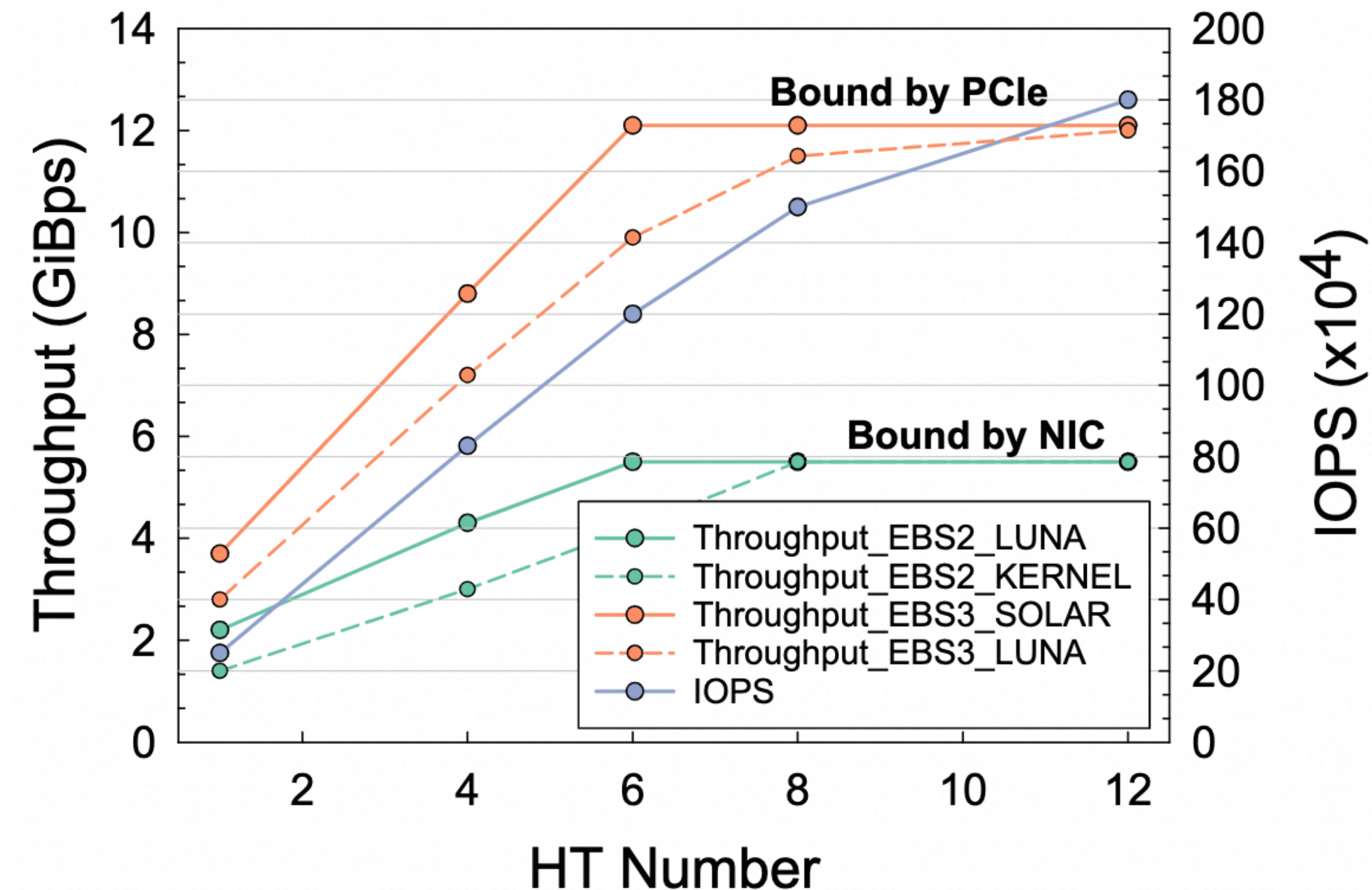
Metrics #2 & #3: IOPS and Throughput

● Upper bound is determined by BlockClient

- ✓ Backend can be easily extended
- ✓ BlockClient is bound by **processing and forwarding** capability
- ✓ From kernel-space to **user-space**, then to **hardware offloading**

● High IOPS/Throughput is often desired but not always needed

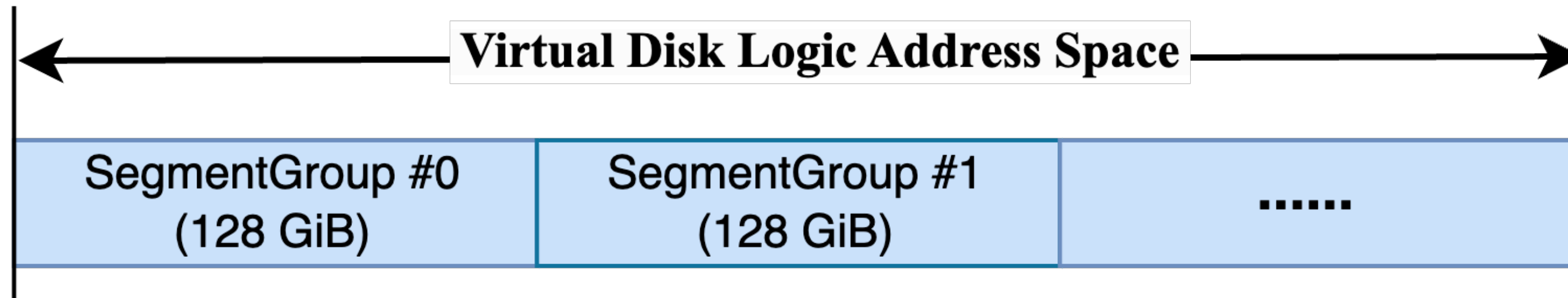
- ✓ Auto performance level (**AutoPL**) Virtual Disk: **on demand without altering the capacity**
- ✓ Base + Burst strategy: **efficiently allocating** IOPS/ throughput to VDs
- ✓ Base throughput means can **definitely** be satisfied
- ✓ Burst throughput means **trying my best** to satisfy



Metrics #4: Capacity

● Flexible space resizing

- ✓ Achieve resizing via adding or removing **SegmentGroups**
- ✓ Virtual disk sizes up to **64 TiB**



● Fast VD cloning

- ✓ *Hard Link* of Pangu files
- ✓ Up to **10,000** virtual disks (each 40 GiB) in **1 min**

Evolving Journey of EBS

Elasticity: A Tale of Four Metrics

Other Topics

● Availability Threats and Solutions *(See Section 4)*

- ✓ **Challenge 1:** a BlockServer crash impacts more VDs
Solution: **Federated BlockManager (Two-layer control nodes)**
- ✓ **Challenge 2:** Segment migration leads to cascading failures
Solution: **Logical Failure Domain (Limited migration)**

● EBS Offloading *(See Section 5)*

- ✓ **FPGA is **not ideal**:** expensive, high failure rates
- ✓ **Blockclient offloading:** **FPGA → ASIC:** 1. cost-friendly 2. a fixed set of functions
- ✓ **BlockServer offloading:** **FPGA → Many-core ARM:** 1. cost-friendly 2. comparable performance

● What if? *(See Section 6)*

- ✓ **Q1: W/o log-structured design?** **Both cost and performance cannot move forward.**
- ✓ **Q2: EBS with open-source software?** **Co-design will be never possible.**
- ✓ **Q3: Not separating Pangu?** **Slow down the development of EBS.**

Thanks

Q & A

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